

Multilevel spectral coarse space methods in FreeFem++ on parallel architectures

Pierre Jolivet

Laboratoire Jacques-Louis Lions
Laboratoire Jean Kuntzmann

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In collaboration with F. Nataf, N. Spillane, V. Dolean,
F. Hecht, C. Prud'homme.

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2 Spectral coarse space

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Construction of a suitable deflation matrix

3 Implementation and numerical results

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Context

Solve large systems arising from the finite element method.

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What are the different alternatives ?

- parallel direct solvers (MUMPS, SuperLU, PaStiX ..),
- parallel iterative solvers (HIPS, Hypre ..),
- hybrid solvers (BoomerAMG, ML, PARDISO ..),
- ...

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⇒ high-performance algorithms on massively parallel distributed memory multiprocessor architectures.

Constraints:

- robust in size and heterogeneities,
- easy to maintain.

Boundary value problems

Given a FE space V and $f \in V^*$, find $u \in V$ such that:

$$a(u, v) = \langle f, v \rangle \quad \forall v \in V \implies Au = F$$

where

$$a_{\Omega}(u, v) = \int_{\Omega} \kappa \nabla u \cdot \nabla v$$

$$a_{\Omega}(u, v) = \int_{\Omega} \mathcal{C} : \varepsilon(u) : \varepsilon(v)$$

$\implies A$ is assumed to be SPD

Notation

Ω is split into N overlapping subdomains $\{\Omega_i\}_{i=1}^N \rightsquigarrow \{V_i\}_{i=1}^N$.

Let:

- R_i be the restriction from V^* to V_i^* (thus R_i^T is the extension by 0 from V_i to V),
- $A_{ij} := R_i A R_j^T$,
- \mathcal{O}_i the set of neighboring subdomains (and $\overline{\mathcal{O}_i} = \mathcal{O}_i \cup \{i\}$),
- $\{\chi_i\}_{i=1}^N$ be a partition of unity:

$$\text{supp}(\chi_i) \subset \Omega_i \quad \text{and} \quad \sum_{i=1}^N \chi_i = 1,$$

- $\{D_i\}_{i=1}^N$ be the discretization of this partition:

$$\sum_{i=1}^N R_i^T D_i R_i = I.$$

One-level Schwarz method

Two classical preconditioners are:

$$M_{\text{RAS}}^{-1} = \sum_{i=1}^N R_i^T D_i A_{ii}^{-1} R_i$$

$$M_{\text{ASM}}^{-1} = \sum_{i=1}^N R_i^T A_{ii}^{-1} R_i$$

Some properties

- obviously parallel,
- not scalable.

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Principle

Let Z be a rectangular matrix so that the “bad eigenvectors” of $M^{-1}A$ belong to the space spanned by its columns. Define:

$$E := Z^T A Z \quad P := I - A Q \quad Q := Z E^{-1} Z^T$$

\implies the number of columns of Z is $O(N)$.

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Two-level preconditioners

$$\mathcal{P}_{\text{BNN}} := (I - A Q)^T M^{-1} (I - A Q) + Q$$

$$\mathcal{P}_{\text{A-DEF1}} := M^{-1} (I - A Q) + Q$$

$$\mathcal{P}_{\text{A-DEF2}} := (I - A Q)^T M^{-1} + Q$$

Differences and similarities studied in (Tang et al. 2009).

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The real problem is: how to build Z ?

A priori construction

Based on an analysis of the underlying PDE, (Nataf et al. 2011; Spillane et al. 2011) — M13P1.

Generalized eigenvalue problem on the overlap

On each Ω_i , find the eigenpairs $(\Lambda_{i_k}, \lambda_{i_k})_k$ such that:

$$a_{\Omega_i}(\Lambda_{i_k}, v_i) = \lambda_{i_k} a_{\Omega_i^\circ}(\chi_i \Lambda_{i_k}, \chi_i v_i) \quad \forall v_i \in V_i$$

which after discretization yields:

$$A_i \Lambda_{i_k} = \lambda_{i_k} D_i A_i^\circ D_i \Lambda_{i_k}$$

A threshold criterion selects only ν_i eigenvectors associated to low frequency eigenvalues:

$$Z = \begin{bmatrix} W_1 & 0 & \cdots & 0 \\ 0 & W_2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & W_N \end{bmatrix}$$

where

$$W_i = \begin{bmatrix} D_i \Lambda_{i_1} & D_i \Lambda_{i_2} & \cdots & D_i \Lambda_{i_{\nu_i}} \end{bmatrix}$$

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Domain Specific Language

FreeFem++, with the help of:

- Metis (Karypis and Kumar 1998)
- SCOTCH (Chevalier and Pellegrini 2008)
- UMFPACK (Davis 2004)
- ARPACK (Lehoucq et al. 1998)
- SLEPc (Hernandez et al. 2005)
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Why use a DS(E)L instead of C/C++/Fortran/... ?

- performances close to low-level language implementation,
- hard to beat something as simple as:

```
varf a(u, v) = int3d(mesh)([dx(u), dy(u), dz(u)]' * [dx(v), dy(v), dz(v)])  
+ int3d(mesh)(f * v) + on(boundary_mesh)(u = 0)
```

Reformulation

Parallel FE computations implies parallel matrix assembly:

$$A_{jk} := R_j A R_k^T \quad \forall k \in \mathcal{O}_j$$

Cannot be done with FreeFem++.

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$$A_{jk} D_k R_k u = R_j R_k^T A_{kk} D_k R_k u$$

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What does it mean ?

$Au =$ black box plugged with a Krylov method

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$$Au = \sum_{j=1}^N AR_j^T D_j R_j u$$

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$$R_i A u = \sum_{j=1}^N R_i A R_j^T D_j R_j u$$

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$$\begin{aligned} R_i A u &= \sum_{j=1}^N R_i A R_j^T D_j R_j u = \sum_{j \in \overline{\mathcal{O}_i}} A_{ij} D_j R_j u \\ &= \sum_{j \in \overline{\mathcal{O}_i}} R_i R_j^T A_{jj} D_j R_j u \end{aligned}$$

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local unknowns on V_j

Coarse space construction

The same goes for the construction of $E := Z^T A Z$.

A double block matrix multiplication

$$E_{ij} = \Lambda_{\mathcal{R}(i)\varphi(i)}^T D_{\mathcal{R}(i)} A_{\mathcal{R}(i)\mathcal{R}(j)} D_{\mathcal{R}(j)} \Lambda_{\mathcal{R}(j)\varphi(j)}$$

where

$$\mathcal{R} : j \mapsto \max \left\{ i : \sum_{k=1}^i \nu_k < j \right\} \quad \varphi : j \mapsto j - \sum_{k=1}^{\mathcal{R}(j)} \nu_k$$

A similar reformulation leads to:

$$E_{ij} = \Lambda_{\mathcal{R}(i)\varphi(i)}^T D_{\mathcal{R}(i)} R_i R_j^T A_{\mathcal{R}(j)\mathcal{R}(j)} D_{\mathcal{R}(j)} \Lambda_{\mathcal{R}(j)\varphi(j)} \neq 0 \iff j \in \mathcal{O}_i$$

Heterogeneous architectures

	N° cores	Memory	Peak performance	Compilers
titane@CEA	12192*	37 To	140 TFLOP/s	Intel
babel@IDRIS	40960	20 To	139 TFLOP/s	IBM + GNU
curie@CEA	80640	322 To	2 PFLOP/s	Intel

* + 46080 CUDA cores

ANR P_{reconditioning} scientific applications on pET_{ascALe} Heterogeneous machines

PRACE High Performance Computing-PDE

[http://www-hpc.cea.fr](http://www-hpc cea.fr), Bruyères-le-Châtel, France.

<http://www.idris.fr>, Orsay, France.

Darcy pressure equation I

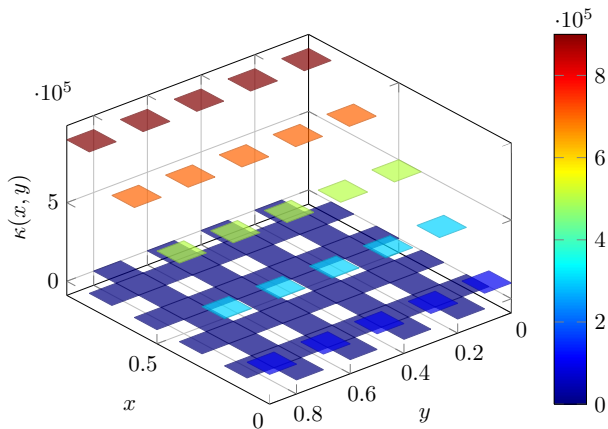


Fig: Two dimensional diffusivity κ

Darcy pressure equation II

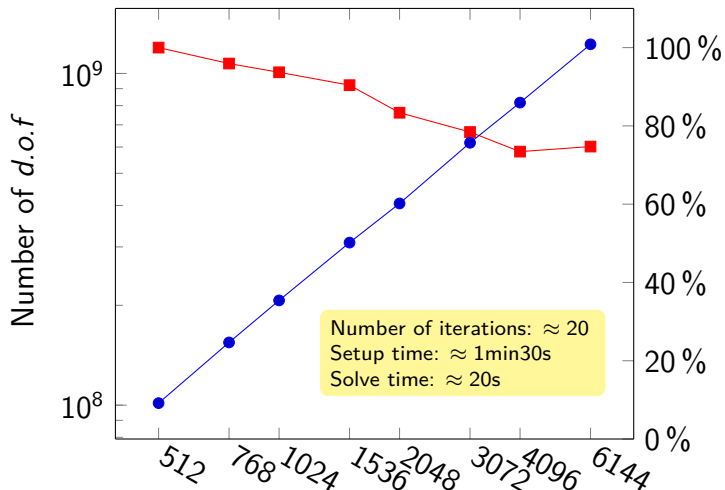


Fig: 2D efficiency (weak scaling) — \mathbb{P}_3 FE, stopping criterion $\varepsilon = 10^{-8}$

System of linear elasticity I (60M unknowns)

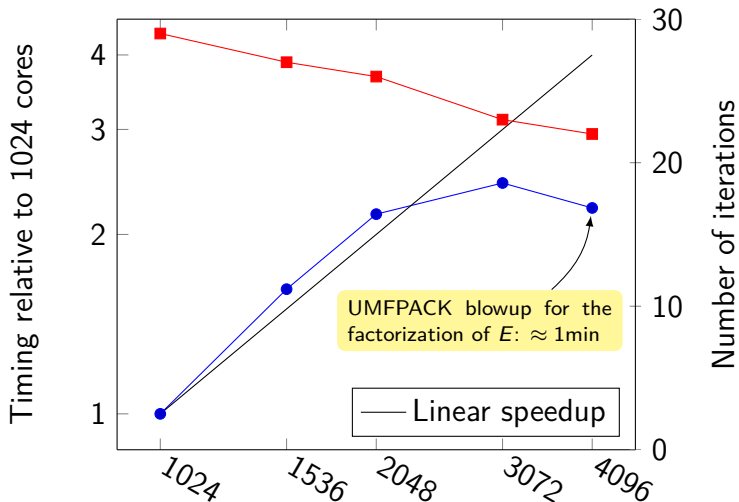


Fig: 3D strong scaling — \mathbb{P}_2 FE, stopping criterion $\varepsilon = 10^{-10}$

System of linear elasticity II (375M unknowns)

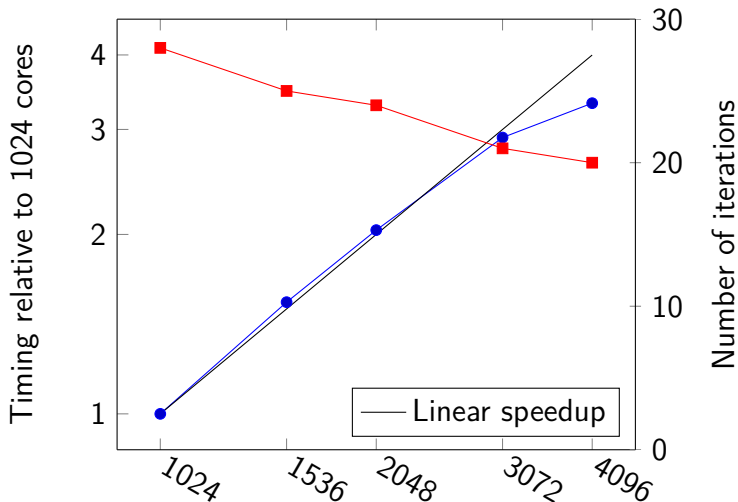


Fig: 2D strong scaling — \mathbb{P}_3 FE, stopping criterion $\varepsilon = 10^{-8}$

The limits reached

6144-way decomposition:

- in \mathbb{R}^2 , 1.2G unkowns,
- in \mathbb{R}^3 , 300M unkowns.

All systems are solved with:

- coarse spaces of size $\llbracket 100; 120\,000 \rrbracket$,
- less than 30 iterations.

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Final words

What to remember:

FreeFem++
+
Two-level DDM = easy framework to solve large systems.

New problems being tackled:

- nonlinear elasticity (Newton-Krylov-Schwarz methods),
- reuse of Krylov and deflation subspaces.

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Thanks for your attention.

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