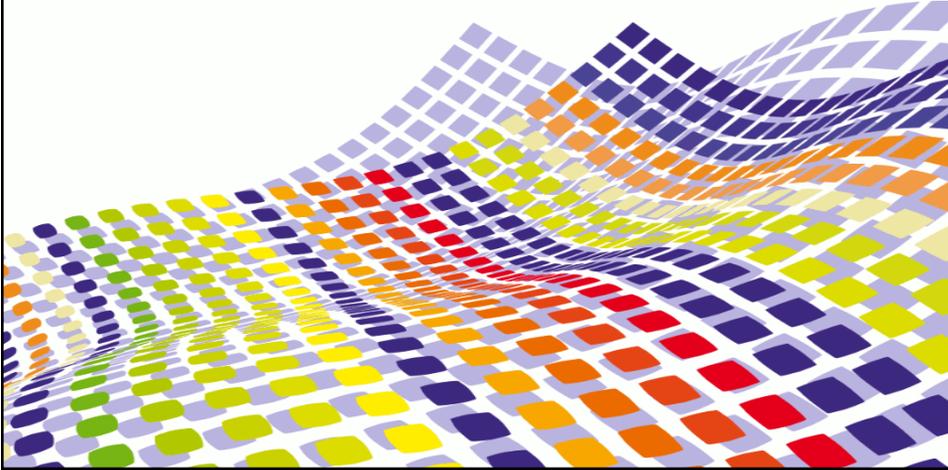




OpenACC Standard

Directives for Accelerators



Credits



- <http://www.openacc.org/>
 - V1.0: November 2011 Specification
- OpenACC, Directives for Accelerators, Nvidia Slideware
- CAPS OpenACC Compiler, HMPP Workbench 3.1.x, CAPS enterprise

Agenda

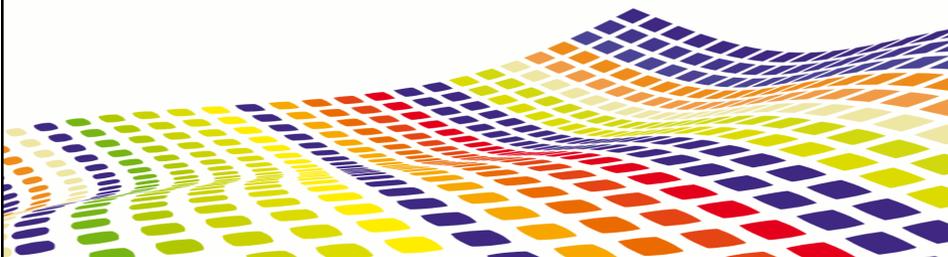


- OpenACC Overview and Compilers
- Programming Model
- Managing Data
- Loops
- Asynchronism
- Runtime API

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OpenACC Overview and Compilers



Directive-based Programming



- Three ways of programming GPGPU applications:

Libraries

Directives

Programming
Languages

*Ready-to-use
Acceleration*

*Quickly Accelerate
Existing Applications*

Maximum Performance

Directive-based Programming



OpenMP

CPU



```
main() {
  double pi = 0.0; long i;

  #pragma omp parallel for reduction(+:pi)
  for (i=0; i<N; i++)
  {
    double t = (double)((i+0.05)/N);
    pi += 4.0/(1.0+t*t);
  }

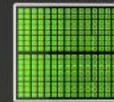
  printf("pi = %f\n", pi/N);
}
```

OpenACC

CPU



GPU



```
main() {
  double pi = 0.0; long i;

  #pragma acc parallel
  for (i=0; i<N; i++)
  {
    double t = (double)((i+0.05)/N);
    pi += 4.0/(1.0+t*t);
  }

  printf("pi = %f\n", pi/N);
}
```

Introduction to Directive-based Programming



- Keeping a unique version of codes, preferably mono-language
 - Reduces maintenance cost
 - Preserves code assets
 - Is less sensitive to fast moving hardware targets
 - Codes last several generations of hardware architecture
- Help to get "portable" performance
 - Multiple forms of parallelism cohabiting
 - Multiple devices (e.g. GPUs) with their own address space
 - Multiple threads inside a device
 - Vector/SIMD parallelism inside a thread
 - Dealing with massive parallelism
- OpenACC is a promising approach

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OpenACC Initiative



- A CAPS, CRAY, Nvidia and PGI initiative
- Open Standard
- A directive-based approach for programming heterogeneous many-core hardware for C and FORTRAN applications
- Available for implementation
 - As CRAY's, PGI's...
 - CAPS OpenACC Compiler → released in April 2012 with HMPP 3.1
 - Satisfies the OpenACC Test Suite provided by University of Houston
- Visit <http://www.openacc-standard.com> for more information

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HMPP Compiler

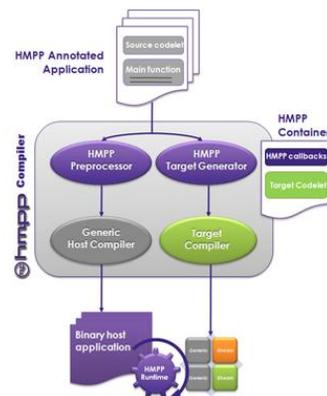


- **The directives**
 - Define hardware implementations of native functions (codelets)
 - Indicate resource allocation and communication
 - Ensure portability (future-proof) and default execution (no exit cost)
- **The toolchain**
 - Helps building manycore applications
 - Includes compilers and target code generators
 - Insulates hardware specific computations
 - Uses hardware vendor SDK
- **The runtime**
 - Helps to adapt to platform configuration
 - Manages hardware resource availability

HMPP Compiler



- **HMPP drives all compilation passes**
 - Host application compilation
 - Calls traditional CPU compilers
 - HMPP Runtime is linked to the host part of the application
 - Device code production
 - According to the specified target
 - A dynamic library is built



```
$ hmpp gcc myprogram.c
$ hmpp gfortran myprogram.f90
```

Programming Model



Execution Model

- Among a bulk of computations executed by the CPU, some regions can be offloaded to hardware accelerators
- Host is responsible for:
 - Allocating memory space on accelerator
 - Initiating data transfers
 - Launching computations
 - Waiting for completion
 - Deallocating memory space
- Accelerators execute parallel regions:
 - Use work-sharing directive
 - Specify level of parallelization

Levels of Parallelism



- Host-controlled execution
- Based on three parallelism levels
 - Gangs – coarse grain
 - Workers – fine grain
 - Vectors – finest grain



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Directive Syntax



- C

```
#pragma acc directive-name [clause [, clause] ...]
{
  code to offload
}
```

- Fortran

```
!$acc directive-name [clause [, clause] ...]
  code to offload
!$acc end directive-name
```

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Work Management: Parallel Construct



- Starts parallel execution on the accelerator
- Creates gangs and workers
- The number of gangs and workers remains constant for the parallel region
- One worker in each gang begins executing the code in the region

```
#pragma acc parallel [...]
{
  ...
}
```

```
$!acc parallel [...]
...
$!acc end parallel
```

Parallel Construct: Gangs and Workers



- The clauses:
 - *num_gangs*
 - *num_workers*

Enables to specify the number of gangs and workers in the corresponding *parallel* section

```
#pragma acc parallel, num_gangs[32], num_workers[256]
{
  ...
  for(i=0; i < n; i++) {
    for(j=0; j < n; j++) {
      ...
    }
  }
  ...
}
```

Work distribution over 32 gangs and 256 workers

Work Management: Kernels Construct



- **Kernels construct**

- Defines a region of code to be compiled into a sequence of accelerator kernels
 - Typically, each loop nest will be a distinct kernel
- The number of gangs and workers can be different for each kernel

```
#pragma acc kernels [...]
{
  for(i=0; i < n; i++) {
    ...
  }
  ...
  for(j=0; j < n; j++) {
    ...
  }
}
```

```

$!acc kernels [...]
  DO i=1,n
  ...
  END DO
  ...
  DO j=1,n
  ...
  END DO
$!acc end kernels

```

1st Kernel

2nd Kernel

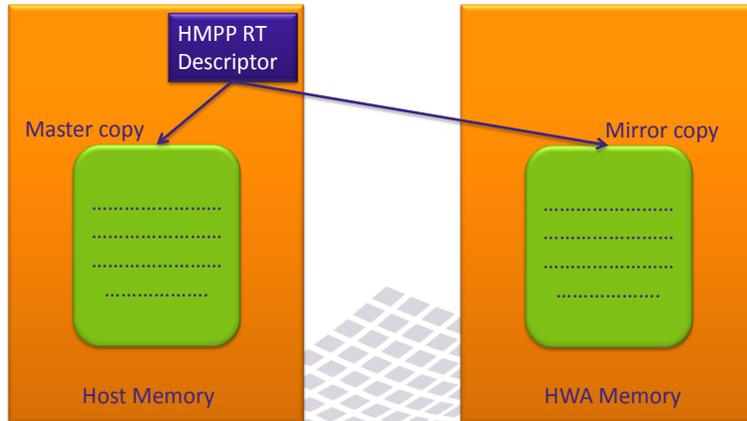
Managing Data



Data Storage



- Mirroring duplicates a CPU memory block into the HWA memory
 - Mirror identifier is a CPU memory block address
 - Only one mirror per CPU block
 - Users ensure consistency of copies via directives



Data Management: Data Constructs



- Defines scalars, arrays and subarrays to be allocated on the device memory for the duration of the region
 - Data can be copied from the host to the device when entering region
 - Data can be copied from the device to the host when exiting region
- *if* clause can be used

```
#pragma acc data [...]
{
  ...
}
```

```
$!acc data [...]
...
$!acc end data
```

Data Allocation: Create Clause



- Declares variable, arrays or subarrays to be allocated in the device memory
- No data specified in this clause will be copied between host and device

```
#pragma acc data, create (A)
{
  ...
}
```

```
$!acc data, create (A)
...
$!acc end data
```

Subarrays



- In C and C++, specified with start and length

```
a[2:n]
```

ie: elements $a[2]$, $a[3]$, ..., $a[2+n-1]$

- If the lower bound is missing, zero is used
- If the length is missing, the difference between the lower bound and the declared size of the array is used

- In Fortran, specified with a list of range specifications

```
a(1:3, 5:6)
```

ie: elements $a(1,5)$, $a(2,5)$, $a(3,5)$, $a(1,6)$, $a(2,6)$, $a(3,6)$

- Any Array or subarray must be a contiguous block of memory

Transfers: Copy Clause



- Declares data that need to be copied from the host to the device when entering the data section
- These data are assigned values on the device that need to be copied back to the host when exiting the data section

```
#pragma acc data, copy (A)
{
  ...
}
```

```
$!acc data, copy (A)
...
$!acc end data
```

Transfers: Copyin/Copyout Clause



- *copyin*
 - Declares data that need to be copied from the host to the device when entering the data section
- *copyout*
 - Declares data that need to be copied from the device to the host when exiting data section

```
#pragma acc data, copyin (A)
{
  ...
}
```

```
$!acc data, copyout (A)
...
$!acc end data
```

Present Clause



- Declares data that are already present on the device
 - Thanks to data region that contains this region of code
- HMPP Runtime will find and use the data on device

```
#pragma acc data, copy (A)
{
  ...
  #pragma acc data, present (A)
  {
    ...
  }
}
```

```
#!acc data, copy (A)
...
#!acc data, present (A)
...
#!acc end data
#!acc end data
```

Data Allocation: Present_or_create Clause



- Declares data that may be present
 - If data is already present, use value in the device memory
 - If not, allocate data on device when entering region and deallocate when exiting
- May be shortened to *pcreate*

```
#pragma acc data, pcreate (A)
{
  ...
}
```

```
#!acc data, pcreate (A)
...
#!acc end data
```

Transfers: Present_or_copy Clause



- If data is already present, use value in the device memory
- If not:
 - Allocates data on device and copies the value from the host at region entry
 - Copies the value from the device to the host and deallocate memory at region exit
- May be shortened to *pcopy*

```
#pragma acc data, pcopy (A)
{
  ...
}
```

```
$!acc data, pcopy (A)
...
$!acc end data
```

Transfers: Present_or_copyin / Present_or_copyout Clause



- If data is already present, use value in the device memory
- If not:
 - Both *present_or_copyin/present_or_copyout* allocate memory on device at region entry
 - *present_or_copyin* copies the value from the host at region entry
 - *present_or_copyout* copies the value from the device to the host at region exit
 - Both *present_or_copyin/present_or_copyout* deallocate memory at region exit
- May be shortened to *pcopyin* and *pcopyout*

```
#pragma acc data, pcopyin (A)
{
  ...
}
```

```
$!acc data, pcopyout (A)
...
$!acc end data
```

Kernels, Parallel Constructs and Data Clauses



- *Kernels* and *parallel* constructs implicitly define data regions
- Data clauses also apply to these structures
- *Kernels* and *parallel* constructs cannot contain other *kernels* or *parallel* regions
- Data inside *kernels* or *parallel* regions data can be managed by a data construct at a higher level

<pre> data.c int A[n] ... #pragma acc data, copyin (A) { ... function (A) ... } </pre>	<pre> kernels.c function(float A[n]) { #pragma acc kernels, \ pcopyin (A) { ... } } </pre>
--	--

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Data Management: Default Behavior



- HMPP compiler is able to detect the variables required on the device for the *kernels* and *parallel* constructs.
- Depending on their type, they follow the following policies
 - Tables: *present_or_copy* behavior
 - Scalar
 - if not live in or live out variable: *private* behavior
 - *copy* behavior otherwise

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Loop Constructs

Kernel Optimization: Loop Construct

- *Loop* directive applies to a loop that immediately follow the directive
- Describes what kind of parallelism to use

```
#pragma acc loop [...]  
for(i=0; i<n; i++)  
{  
  ...  
}
```

```
$(acc loop [...]  
DO i=1,n  
  ...  
END DO
```

Sequential Execution



- The *seq* clause specifies that the associated loop should be executed sequentially
- This is the default behavior in a *parallel* region

```
#pragma acc loop seq
for(i=0; i<n; i++)
{
  ...
}
```

```
!acc loop seq
DO i=1,n
  ...
END DO
```

Data Independence



- The clause *independent* specifies that iterations of the loop are data-independent
- Allowed on loop directives in kernels regions
- Allows the compiler to generate code to execute the iterations in parallel with no synchronisation

```
#pragma acc loop independent
for(i=0; i<n; i++)
{
  for(j=0; j<m; j++)
  {
    A(j,i*3+MOD(i,2)) = i*j;
  }
}
```

```
#pragma acc loop independent
for(i=0; i<n; i++)
{
  for(j=0; j<m; j++)
  {
    A(j,i*3+MOD(i,2)) =
      i*A(i,j-1);
  }
}
```

Programming error

Gangs



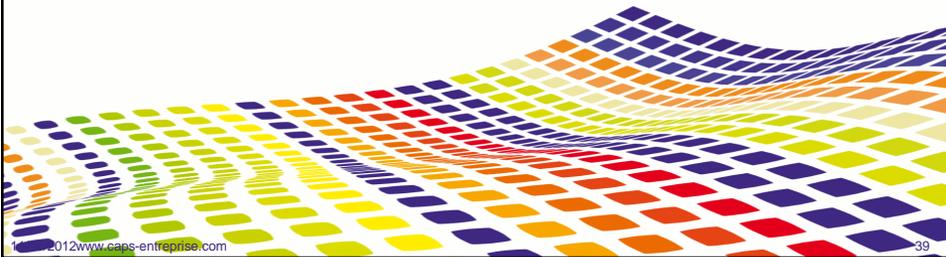
- **Gang clause:**
 - The iterations of the following loop are executed in parallel
 - In a parallel construct:
 - Iterations are distributed among the gangs created by the *parallel* construct
 - No argument is allowed
 - In a kernels construct
 - Iterations are distributed among the gangs created by the kernel created by a loop
 - An argument can specify the number of gangs to use for this loop

Workers



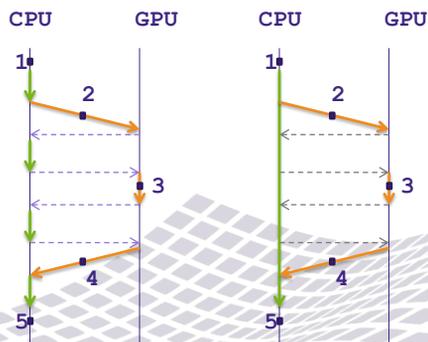
- **Worker clause:**
 - The iterations of the following loop are executed in parallel
 - In a parallel construct:
 - Iterations are distributed among the multiple workers within a single gang
 - No argument is allowed
 - Loop iterations must be data independent, unless it performs a reduction operation
 - In a kernels construct:
 - Iterations are distributed among the workers within the gangs created by the kernel within a loop
 - An argument can specify the number of workers to use for this loop

Asynchronism



Asynchronism

- By default, the code on the accelerator is synchronous
 - The host waits for completion of the parallel or kernels region
- The *async* clause enables to use the device while the host process continues with the code following the region
- Can be used on *parallel* and *kernels* regions and *update* directives



Wait Directive



- Causes the program to wait for an asynchronous activity
 - Parallel, kernels regions or update directives
- An identifier can be added to the async clause and wait directive:
 - Host thread will wait for the asynchronous activities with the same ID
- Without any identifier, the host process waits for all asynchronous activities

```
#pragma acc kernels, async
{
  ...
}
#pragma acc kernels, async
{
  ...
}
#pragma acc wait
```

```
$!acc kernels, async 1
...
$!acc end kernels
...
$!acc kernels, async 2
...
$!acc end kernels
...
$!acc wait 1
```

Runtime API



Runtime Library Definition



- For C:
 - Header file: openacc.h
- For Fortran:
 - Interface declaration in: openacc_lib.h in a Fortran module called openacc
- **acc_device_t: type of accelerator device**
 - acc_device_none
 - acc_device_default
 - acc_device_host
 - acc_device_not_host
 - ...

Runtime API



- **int acc_get_num_device (acc_device_t) (C)**
- **integer function acc_get_num_device (devicetype) (Fortran)**
 - Returns the number of accelerator devices of the given type attached to the host
- **int acc_set_device_type (acc_device_t) (C)**
- **subroutine acc_set_device_type (devicetype) (Fortran)**
 - Tells the runtime which type of device to use
- **acc_device_type acc_get_device_type (void) (C)**
- **function acc_get_device_type () (Fortran)**
 - Tells the program what type of device will be used

Runtime API



- `void acc_set_device_num (int, acc_device_t) (C)`
- subroutine `acc_set_device_num (devicenum, devicetype) (Fortran)`
 - Tells the runtime which device to use
- `int acc_get_device_num (acc_device_t) (C)`
- Integer function `acc_get_device_num (devicetype) (Fortran)`
 - Return the device number of the specified device type that will be used

Runtime API



- `void acc_init (acc_device_t) (C)`
- Subroutine `acc_init (devicetype) (Fortran)`
 - Initialize the runtime for the given type
- `void acc_shutdown (acc_device_t) (C)`
- Subroutine `acc_shutdown (devicetype) (Fortran)`
 - Disconnect the program from the accelerator device
- `void* acc_malloc (size_t) (C)`
 - Allocates memory on accelerator device
 - Pointers assigned to this function may be reused
- `void* acc_free (size_t) (C)`
 - Deallocates memory on accelerator device

Conclusion



- Beware of compiler-dependent behaviors
- Fast development of high-level heterogeneous applications
 - For C and FORTRAN code
- Explicit the calls to a hardware accelerator in your code
 - Whatever the target
 - CAPS OpenACC compiler supports:
 - Nvidia Tesla GPUs
 - AMD
 - X86 Intel Phi

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CAPS

Accelerator Programming Model Parallelization

Directive-based programming GPGPU **Manycore programming**

Hybrid Manycore Programming HPC community OpenACC

Petaflops Parallel computing HPC open standard

Multicore programming Exaflops NVIDIA Cuda

Code speedup Hardware accelerators programming

High Performance Computing OpenHMPP

Parallel programming interface

Massively parallel

Open CL

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