Analysis and Optimization of the Memory Access Behavior of Applications

Ecole
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My Background

• Chair for computer architecture at CS faculty, TUM
  – how to exploit current & future (HPC) systems (multicore, accelerators)
  – programming models, performance analysis tools, application tuning
• PhD on load balancing of commercial car crash code (MPI) 2003
• Interested especially in cache analysis and optimization
  – cache simulation: Callgrind (using Valgrind)
  – applied to 2D/3D stencil codes
  – recently extended to multicore (new bottlenecks, new benefits)
• Invited by Romaric David to give a talk at this workshop
Topic of this Morning: Bottleneck Memory

- Why should you care about memory performance?
- Most (HPC) applications often do memory accesses
- Good vs. bad use of the memory hierarchy can be ~ factor 100 (!)
- Example: modern processor with 3GHz clock rate, 2 sockets
  - latency to remote socket ~ 100 ns: 300 clock ticks
  - bandwidth (1 core) ~ 15 GB/s
  - compare to L1 access: latency 2-3 ticks, bandwidth ~150GB/s
- Bad memory performance easily can dominate performance (better memory performance also will speed up parallel code)
Topic of this Morning: Bottleneck Memory

Still getting more important
• compute power on one chip still increases
• main memory latency will stay (off-chip distance)
• bandwidth increases, but not as much as compute power
⇒ Memory Wall (stated already in 1994)

In addition:
• with multi-core, cores share connection to main memory!
The Memory Wall

CPU Peak Performance (clock & cores)
+ 40% / year

Main Memory Performance
+ 7% / year

Access latency to main memory today up to 300 cycles
Assume 2 Flops/clock ticks ➔ 600 Flops wasted while waiting for one main memory access!
Topic of this Morning: Bottleneck Memory

- Further getting more important not only for performance, but for problem no.1 in the future: power consumption (Power Wall)
  - reason that we have multi-core today
  - most significant cost factor for compute centers in the future
  - users not to be charged by hours, but by energy consumption?

- Comparison computation vs. memory access [Dongarra, PPAM 2011]

- today: for 1 memory access saved, can do 48 FMAs more
  2018: 192 FMAs more

- solution (?): do redundant calculation to avoid memory access
Outline: Part 1

The Memory Hierarchy
   Caches: Why & How do they work?

Bad Memory Access Patterns
   How to not exploit Caches

Cache Optimization Strategies
   How to exploit Caches even better
Outline: Part 2

Cache Analysis
   Measuring on real Hardware vs. Simulation

Cache Analysis Tools

Case Studies

Hands-on
The Memory Hierarchy

Two facts of modern computer systems
• processor cores are quite fast
• main memory is quite slow

Why? Different design goals
• everybody wants a fast processor
• everybody wants large amounts of cheap memory

Why is this not a contradiction? There is a solution to bridge the gap:
• a hierarchy of buffers between processor and main memory
• often effective, and gives seemingly fast and large memory
Solution: The Memory Hierarchy

We can build very fast memory (for a processor), but

- it has to be small (only small number of cascading gates)
  - tradeoff: buffer size vs. buffer speed
- it has to be near (where data is to be used)
  - on-chip, not much space around execution units
- it will be quite expensive (for its size)
  - SRAM needs a lot more energy and space than DRAM

➤ use fast memory only for data most relevant to performance
➤ if less relevant, we can afford slower access, allowing more space
➤ this works especially well if “most relevant data” fits into fast buffer
Solution: The Memory Hierarchy

<table>
<thead>
<tr>
<th>Size</th>
<th>Latency</th>
<th>Bandwidth</th>
</tr>
</thead>
<tbody>
<tr>
<td>300 B</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>32 kB</td>
<td>3</td>
<td>100 GB/s</td>
</tr>
<tr>
<td>4 MB</td>
<td>20</td>
<td>30 GB/s</td>
</tr>
<tr>
<td>4 GB</td>
<td>200</td>
<td>15 GB/s</td>
</tr>
<tr>
<td>4 GB</td>
<td>300</td>
<td>10 GB/s</td>
</tr>
<tr>
<td>1 TB</td>
<td>&gt; 10^7</td>
<td>0.2 GB/s</td>
</tr>
</tbody>
</table>
Solution: The Memory Hierarchy

Programmers want memory to be a flat space
• registers not visible, used by compilers
• on-chip buffers are
  – not explicitly accessed, but automatically filled from lower levels
  – indexed by main memory address
  – hold copies of blocks of main memory
 ➔ not visible to programmers: caches
• transparent remote memory access provided by hardware
• extension on I/O devices by MMU & OS

Let’s concentrate on Processor Caches…
Solution: Processor Caches

Why are Caches effective? Because typical programs
• often access same memory cells repeatedly
  – temporal locality \(\Rightarrow\) good to keep recent accessed data in cache
• often access memory cells near recent accesses
  – spatial locality \(\Rightarrow\) good to work on blocks of nearside data (cache line)

“Principle of Locality”

So what’s about the Memory Wall?
• the degree of “locality” depends on the application
• at same locality, the widening gap between processor and memory performance reduces cache effectiveness
Example: Sequence with 10 Accesses

- memory latency: 3
- cache latency: 1
- without cache: 30
- cache exploiting temporal locality: 22 (6 misses, 4 hits)
- cache exploiting temporal and spatial locality: 16 (3 misses, 7 hits)
Basic Cache Properties (1)

- Cache holds copies of memory blocks
  - space for one copy called “cache line” \(\rightarrow\) **Cache Line Size**
  - transfers from/to main memory always at line size granularity
- Cache has restricted size: **Cache Size**
  - line size 2, cache size 6 (= 3 lines)
  - line size 2, cache size 4 (=2 lines)
- Which copy to evict for new copy
  - **Replacement Policy**
  - Typically: Evict Least Recently Used (LRU)
Basic Cache Properties (2)

- every cache line knows the memory address it has a copy of („tag“)
- comparing all tags at every access ➞ expensive (space & energy)
- better: reduce number of comparisons per access
  - group cache lines into sets
  - a given address can only be stored into a given set
  - lines per set: **Associativity**
- example: 2 lines ( ), sequence 1/3/1/3/2/4/2/4
Solution: Processor Caches

The “Principle of Locality” makes caches effective
- How to improve on that?
- Try to further reduce misses!

Options
- increase cache line size!
  - can reduce cache effectiveness, if not all bytes are accessed
- predict future accesses (hardware prefetcher), load before use
  - example: stride detectors (more effective if keyed by instruction)
  - allows “burst accesses” with higher netto bandwidth
  - only works if bandwidth not exploited anyway (demand vs. speculative)
  - can increase misses if prefetching is too aggressive
The Memory Hierarchy on Multi-Core

Principle of Locality often holds true across multiple threads
• example: threads need same vectors/matrices
• caches shared among cores can be beneficial
• sharing allows threads to prefetch data for each other

However, if threads work on different data…
• example: disjunct partitioning of data among threads
• threads compete for space, evict data of each other
• trade-off: only use cache sharing on largest on-chip buffer
The Memory Hierarchy on Multi-Core

Typical example (modern Intel / AMD processors)

Why are there 3 levels?
• cache sharing increases on-chip bandwidth demands by cores
• L1 is very small to be very fast ➔ still lots of references to L2
• private L2 caches reduce bandwidth demands for shared L3
Caches and Multi-Processor Systems

The Cache Coherence Problem

• suppose 2 processors/cores with private caches at same level
• P1 reads a memory block X
• P2 writes to the block X
• P1 again reads from block X (which now is invalid!)

A strategy is needed to keep caches coherent

• writing to X by P2 needs to invalidate or update copy of X in P1
• cache coherence protocol
• all current multi-socket/-core systems have fully automatic cache coherence in hardware (today already a significant overhead!)
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   How to not exploit Caches

Cache Optimization Strategies
   How to exploit Caches even better
Memory Access Behavior

How to characterize good memory access behavior?

**Cache Hit Ratio**

- percentage of accesses which was served by the cache
- good ratio: > 97%

Symptoms of bad memory access: Cache Misses

Let’s assume that we can not change the hardware as countermeasure for cache misses (e.g. enlarging cache size)
Memory Access Behavior: Cache Misses

Classification:

• **cold / compulsory** miss
  – first time a memory block was accessed

• **capacity** miss
  – recent copy was evicted because of too small cache size

• **conflict** miss
  – recent copy was evicted because of too low associativity

• **concurrency** miss
  – recent copy was evicted because of invalidation by cache coherence protocol

• **prefetch inaccuracy** miss
  – recent copy was evicted because of aggressive/imprecise prefetching
Bad Memory Access Behavior (1)

Lots of cold misses

- Each memory block only accessed once, and
- Prefetching not effective because accesses are not predictable or bandwidth is fully used
- Usually not important, as programs access data multiple times
- Can become relevant if there are lots of context switches (when multiple processes synchronize very often)
  - L1 gets flushed because virtual addresses get invalid
Bad Memory Access Behavior (2)

Lots of capacity misses
• blocks are only accessed again after eviction due to limited size
  – number of other blocks accessed in-between (= reuse distance) > number of cache lines
  – example: sequential access to data structure larger than cache size
• and prefetching not effective

Countermeasures
• reduce reuse distance of accesses = increase temporal locality
• improve utilization inside cache lines = increase spatial locality
• do not share cache among threads accessing different data
• increase predictability of memory accesses
Bad Memory Access Behavior (3)

Lots of conflict misses

- blocks are only accessed again after eviction due to limited set size
- example:
  - matrix where same column in multiple rows map to same set
  - we do a column-wise sweep

![Diagram showing memory blocks assigned to different sets.](image)
Bad Memory Access Behavior (3)

Lots of conflict misses
• blocks are only accessed again after eviction due to limited set size

Countermeasures
• set sizes are similar to cache sizes: see last slide…
• make successive accesses cross multiple sets
Bad Memory Access Behavior (4)

Lots of concurrency misses
• lots of conflicting accesses to same memory blocks by multiple processors/cores, which use private caches
  – “conflicting access”: at least one processor is writing

Two variants: same block is used
• because processors access same data
• even though different data are accessed, the data resides in same block (= false sharing)
  – example: threads often write to nearside data
    (e.g. using OpenMP dynamic scheduling)
Bad Memory Access Behavior (4)

Lots of concurrency misses

- lots of conflicting accesses to same memory blocks by multiple processors/cores, which use private caches

Countermeasures

- reduce frequency of accesses to same block by multiple threads
- move data structures such that data accessed by different threads reside on their own cache lines
- place threads to use a shared cache
Bad Memory Access Behavior (5)

Lots of prefetch inaccuracy misses

- much useful data gets evicted due to misleading access patterns
- example: prefetchers typically “detect” stride pattern after 3-5 regular accesses, prefetching with distance 3-5
  - frequent sequential accesses to very small ranges (5-10 elements) of data structures

Countermeasures

- use longer access sequences with strides
- change data structure if an access sequence accidently looks like a stride access
Memory Access Behavior: Cache Misses

Classifications:
- kind of misses
- each cache miss needs another line to be evicted: is the previous line modified (= dirty) or not?
  - yes: needs write-back to memory
  - increases memory access latency
Outline: Part 1

The Memory Hierarchy
- Caches: Why & How do they work?

Bad Memory Access Patterns
- How to not exploit Caches

Cache Optimization Strategies
- How to exploit Caches even better
The Principle of Locality is not enough...

Reasons for Performance Loss for SPEC2000
[Beyls/Hollander, ICCS 2004]
Basic efficiency guidelines

Always use a performance analysis tool before doing optimizations: How much time is wasted where because of cache misses?

1. Choose the best algorithm
2. Use efficient libraries
3. Find good compiler and options (“-O3”, “-fno-alias” …)
4. Reorder memory accesses
5. Use suitable data layout
6. Prefetch data

Cache Optimizations

Warning: Conflict and capacity misses are not easy to distinguish…
Cache Optimization Strategies: Reordering Accesses

- **Blocking**: make arrays fit into a cache
Cache Optimization Strategies: Reordering Accesses

- Blocking: make arrays fit into a cache
- Blocking in multiple dimensions (example: 2D)
Cache Optimization Strategies: Reordering Accesses

- Blocking: make arrays fit into a cache
- Blocking in multiple dimensions (example: 2D)
- Nested blocking: tune to multiple cache levels
  - can be done recursively according to a space filling curve
  - example: Morton curve (without “jumps”: Hilbert, Peano…)
  - cache-oblivious orderings/algorithms (= automatically fit to varying levels and sizes using the same code)

Cache Optimization Strategies: Reordering Accesses

- Extreme blocking with size 1: Interweaving
  - combined with blocking in other dimensions, results in pipeline patterns
  - On multi-core: consecutive iterations on cores with shared cache

- Block Skewing:
  Change traversal order over non-rectangular shapes

- For all reorderings: preserve data dependencies of algorithm!
Cache Optimization Strategies: Suitable Data Layout

Strive for best spatial locality

• use compact data structures
  (arrays are almost always better than linked lists!)
• data accessed at the same time should be packed together
• avoid putting frequent and rarely used data packed together
• object-oriented programming
  – try to avoid indirections
  – bad: frequent access of only one field of a huge number of objects
  – use proxy objects, and structs of arrays instead of arrays of structs
• best layout can change between different program phases
  – do format conversion if accesses can become more cache friendly
  – (also can be important to allow for vectorization)
Cache Optimization Strategies: Prefetching

Allow hardware prefetcher to help loading data as much as possible
• make sequence of memory accesses predictable
  – prefetchers can detect multiple streams at the same time (>10)
• arrange your data accordingly in memory
• avoid non-predictable, random access sequences
  – pointer-based data structures without control on allocation of nodes
  – hash tables accesses

Software controlled prefetching (difficult !)
• switch between block prefetching & computation phases
• do prefetching in another thread / core („helper thread“)
Countermeasures for Capacity Misses

Reduce reuse distance of accesses = increase temporal locality

Strategy:
• blocking

Effectiveness can be seen by
• reduced number of misses
• in reuse distance histogram (needs cache simulator)
Countermeasures for Capacity Misses

Improve utilization inside cache lines = increase spatial locality

Strategy:
• improve data layout

Effectiveness can be seen by
• reduced number of misses
• spatial loss metric (needs cache simulator)
  – counts number of bytes fetched to a given cache level but never actually used before evicted again
• spatial access homogeneity (needs cache simulator)
  – variance among number of accesses to bytes inside of a cache line
Countermeasures for Capacity Misses

Do not share cache among threads accessing different data

Strategy:
- explicitly assign threads to cores
- “sched_setaffinity” (automatic system-level tool: autopin)

Effectiveness can be seen by
- reduced number of misses
Countermeasures for Capacity Misses

Increase predictability of memory accesses

Strategy:
• improve data layout
• reorder accesses

Effectiveness can be seen by
• reduced number of misses
• performance counter for hardware prefetcher
• run cache simulation with/without prefetcher simulation
Countermeasures for Conflict Misses

Make successive accesses cross multiple cache sets

Strategy:
• change data layout by Padding
• reorder accesses

Effectiveness can be seen by
• reduced number of misses
Countermeasures for Concurrency Misses

Reduce frequency of accesses to same block by multiple threads

Strategy:

• for true data sharing: do reductions by partial results per thread
• for false sharing (reduce frequency to zero = data accessed by different threads reside on their own cache lines)
  – change data layout by padding (always possible)
  – change scheduling (e.g. increase OpenMP chunk size)

Effectiveness can be seen by

• reduced number of concurrency misses (there is a perf. counter)
Countermeasures for Misses triggering Write-Back

Only general rule:
• Try to avoid writing if not needed

Sieve of Eratosthenes:  
\begin{verbatim}
isPrim[*] = 1;
for(i=2; i<n/2; i++)
   if (isPrim[i] == 1)
      for(j=2*i; i<n; j+=i)
         isPrim[j] = 0;
\end{verbatim}

~ 2x faster (!):  
\begin{verbatim}
isPrim[*] = 1;
for(i=2; i<n/2; i++)
   if (isPrim[i] == 1)
      for(j=2*i; i<n; j+=i)
         if (isPrim[j]==1)
            isPrim[j] = 0;
\end{verbatim}
Outline: Part 2

Cache Analysis
  Measuring on real Hardware vs. Simulation

Cache Analysis Tools

Case Studies

Hands-on
Sequential Performance Analysis Tools

Count occurrences of events

• resource exploitation is related to events
• SW-related: function call, OS scheduling, ...
• HW-related: FLOP executed, memory access, cache miss, time spent for an activity (like running an instruction)

Relate events to source code

• find code regions where most time is spent
• check for improvement after changes
• „Profile“: histogram of events happening at given code positions
• inclusive vs. exclusive cost
How to measure Events (1)

Where?

• on real hardware
  – needs sensors for interesting events
  – for low overhead: hardware support for event counting
  – difficult to understand because of unknown micro-architecture, overlapping and asynchronous execution

• using machine model
  – events generated by a simulation of a (simplified) hardware model
  – no measurement overhead: allows for sophisticated online processing
  – simple models relatively easy to understand

Both methods have pro & contra, but reality matters in the end
How to measure Events (2)

SW-related
- instrumentation (= insertion of measurement code)
  - into OS / application, manual/automatic, on source/binary level
  - on real HW: always incurs overhead which is difficult to estimate

HW-related
- read Hardware Performance Counters
  - gives exact event counts for code ranges
  - needs instrumentation
- statistical: **Sampling**
  - event distribution over code approximated by every N-th event
  - HW notifies only about every N-th event \(\Rightarrow\) Influence tunable by N
Outline: Part 2

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Hands-on
Analysis Tools

- **GProf**
  - Instrumentation by compiler for call relationships & call counts
  - Statistical time sampling using timers
  - Pro: available almost everywhere (gcc: -pg)
  - Contra: recompilation, measurement overhead, heuristic

- **Intel VTune (Sampling mode) / Linux Perf (>2.6.31)**
  - Sampling using hardware performance counters, no instrumentation
  - Pro: minimal overhead, detailed counter analysis possible
  - Contra: call relationship can not be collected
    (this is not about call stack sampling: provides better context…)

- **Callgrind**: machine model simulation
Callgrind: Basic Features

Based on Valgrind

- runtime instrumentation infrastructure (no recompilation needed)
- dynamic binary translation of user-level processes
- Linux/AIX/OS X on x86, x86-64, PPC32/64, ARM

- correctness checking & profiling tools on top
  - “memcheck”: accessibility/validity of memory accesses
  - “helgrind” / ”drd”: race detection on multithreaded code
  - “cache-grind”/”callgrind”: cache & branch prediction simulation
  - “massif”: memory profiling

- Open source (GPL), www.valgrind.org
Callgrind: Basic Features

Measurement

- profiling via machine simulation (simple cache model)
- instruments memory accesses to feed cache simulator
- hook into call/return instructions, thread switches, signal handlers
- instruments (conditional) jumps for CFG inside of functions

Presentation of results

- callgrind_annotate
- \{Q,K\}Cachegrind
Pro & Contra (i.e. Simulation vs. Real Measurement)

Usage of Valgrind
- driven only by user-level instructions of one process
- slowdown (call-graph tracing: 15-20x, + cache simulation: 40-60x)
  - “fast-forward mode”: 2-3x
  - allows detailed (mostly reproducable) observation
  - does not need root access / can not crash machine

Cache model
- “not reality”: synchronous 2-level inclusive cache hierarchy
  (size/associativity taken from real machine, always including LLC)
  - easy to understand / reconstruct for user
  - reproducible results independent on real machine load
  - derived optimizations applicable for most architectures
Callgrind: Usage

- `valgrind -tool=callgrind [callgrind options] yourprogram args`
- `cache simulator: --cache-sim=yes`
- `branch prediction simulation (since VG 3.6): --branch-sim=yes`
- `enable for machine code annotation: --dump-instr=yes`
- `start in “fast-forward”: --instr-atstart=yes`
  - switch on event collection: `callgrind_control -i on/ Macro`
- `spontaneous dump: callgrind_control -d [dump identification]`
- `current backtrace of threads (interactive): callgrind_control -b`
- `separate dumps per thread: --separate-threads=yes`
- `cache line utilization: --cacheuse=yes`
- `enable prefetcher simulation: --simulate-hwpref=yes`
- `jump-tracing in functions (CFG): --collect-jumps=yes`
KCacheGrind: Features

- open source, GPL, kcachegrind.sf.net
- included with KDE3 & KDE4

Visualization of
- call relationship of functions (callers, callees, call graph)
- exclusive/Inclusive cost metrics of functions
  - grouping according to ELF object / source file / C++ class
- source/assembly annotation: costs + CFG
- arbitrary events counts + specification of derived events

Callgrind support (file format, events of cache model)
KCacheGrind: Usage

\{k,q\} cachegrind callgrind.out.<pid>

- left: “Dockables”
  - list of function groups according to
    - library (ELF object)
    - source
    - class (C++)
  - list of functions with
    - inclusive
    - exclusive costs

- right: visualization panes
Visualization panes for selected function

- List of event types
- List of callers/callees
- Treemap visualization
- Call Graph
- Source annotation
- Assemly annotation
Weidendorfer: Memory Access Analysis and Optimization
Outline: Part 2

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Hands-on
Case Studies

- Get ready for hands-on
  - matrix multiplication
  - 2D relaxation
Matrix Multiplication

- Kernel for $C = A \times B$
  - Side length $N \Rightarrow N^3$ multiplications + $N^3$ additions

\[
c[k][i] = a[k][j] \times b[j][i]
\]
Matrix Multiplication

- Kernel for $C = A \times B$
  - 3 nested loops $(i,j,k)$: What is the best index order? Why?

  ```
  for (i=0; i<N; i++)
    for (j=0; j<N; j++)
      for (k=0; k<N; k++)
        c[k][i] = a[k][j] * b[j][i]
  ```

  - blocking for all 3 indexes, block size $B$, $N$ multiple of $B$

  ```
  for (i=0; i<N; i+=B)
    for (j=0; j<N; j+=B)
      for (k=0; k<N; k+=B)
        for (ii=i; ii<i+B; ii++)
          for (jj=j; jj<j+B; jj++)
            for (kk=k; kk<k+B; kk++)
              c[k+kk][i+ii] = a[k+kk][j+jj] * b[j+jj][i+ii]
  ```
Iterative Solver for PDEs: 2D Jacobi Relaxation

Example: Poisson

One iteration:

```c
for (i=1; i<N-1; i++)
    for (j=1; j<N-1; j++)
        u2[i][j] = ( u[i-1][j] + u[i][j-1] + u[i+1][j] + u[i][j+1] ) / 4.0;

u[*][*] = u2[*][*];
```

Optimization: Interleave 2 iterations

- iteration 1 for row 1
- iteration 1 for row 2, iteration 2 for row 1
- iteration 1 for row 3, iteration 2 for row 2
- ...
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Hands-on
How to run with MPI

• Run valgrind with mpirun (bt-mz: example from NAS)
  
  ```sh
  export OMP_NUM_THREADS=4
  mpirun -np 4 valgrind --tool=callgrind --cache-sim=yes \ 
  --separate-threads=yes ./bt-mz_B.4
  ```

• load all profile dumps at once:
  
  – run in new directory, “qcachegrind callgrind.out”
Getting started / Matrix Multiplication / Jacobi

• Try it out yourself (on intelnode)
  “cp -r /srv/app/kcachegrind/kcg-examples .”
  example exercises are in “exercises.txt”

• What happens in „/bin/ls“?
  – valgrind --tool=callgrind ls /usr/bin
  – qcachegrind
  – What function takes most instruction executions? Purpose?
  – Where is the main function?
  – Now run with cache simulation: --cache-sim=yes